

P1

IS THERE A SPIN GLASS IN
YOUR CELLULAR PHONE ?

YES, AND IT'S TRYING TO
MINIMIZE ITS BETHE FREE-ENERGY.

[OR THERE WILL BE SOON!!]

A.M.

LPTENS, PARIS

N.SOURLAS

S.FRANZ

M.LEONE

F.RICCI-TERSENGHI

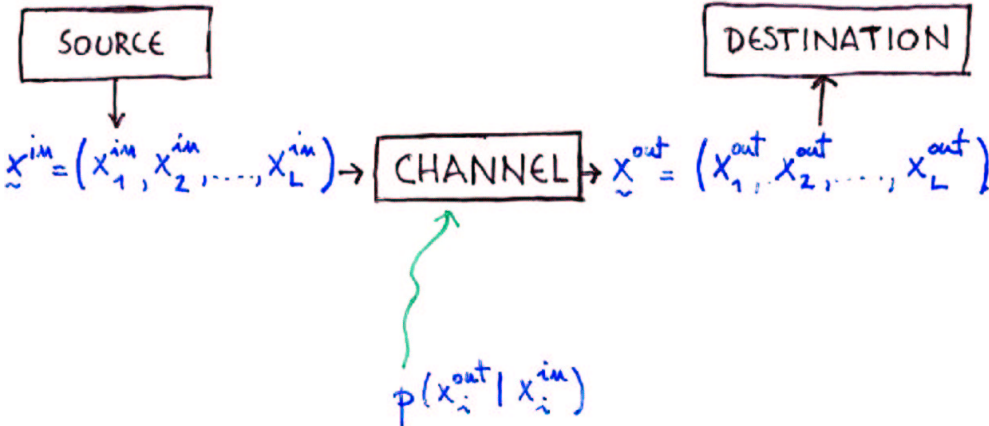
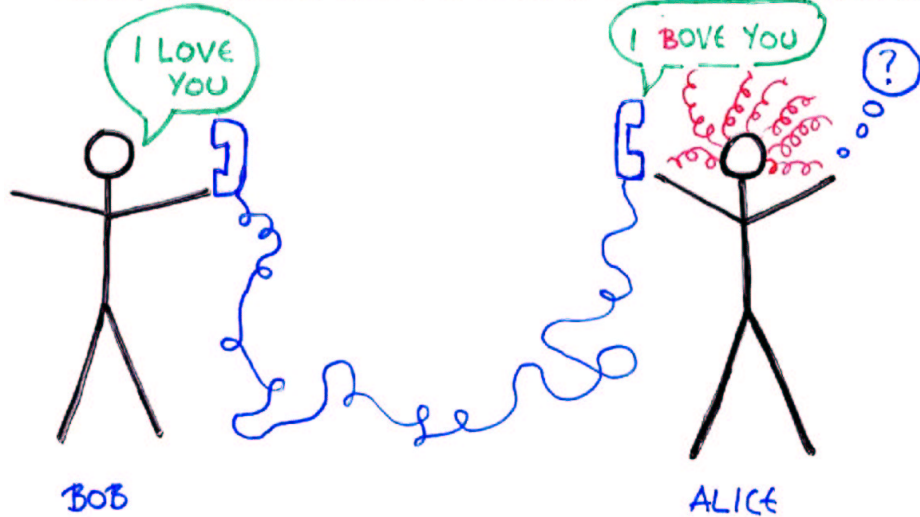
P2

PLAN :

- INTRODUCTION : SHANNON THEORY
- INTRODUCTION : LOW DENSITY PARITY CHECK CODES
- INTRODUCTION : ITERATIVE DECODING
- THE PHASE DIAGRAM
[OR: WHAT CODING THEORY CAN LEARN FROM PHYSICS]
- CONCLUSION
[OR : WHAT PHYSICS CAN LEARN FROM CODING THEORY]

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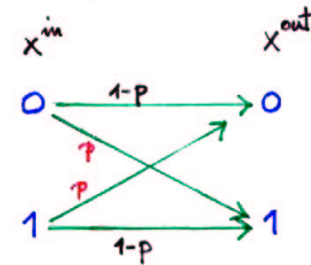
INTRODUCTION : SHANNON THEORY



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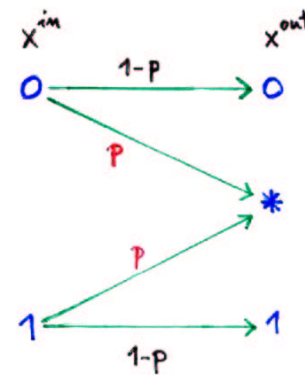
NOISY CHANNELS

- example 1 : BINARY SYMMETRIC CHANNEL (BSC)



$p \equiv$ "FLIP PROBABILITY"

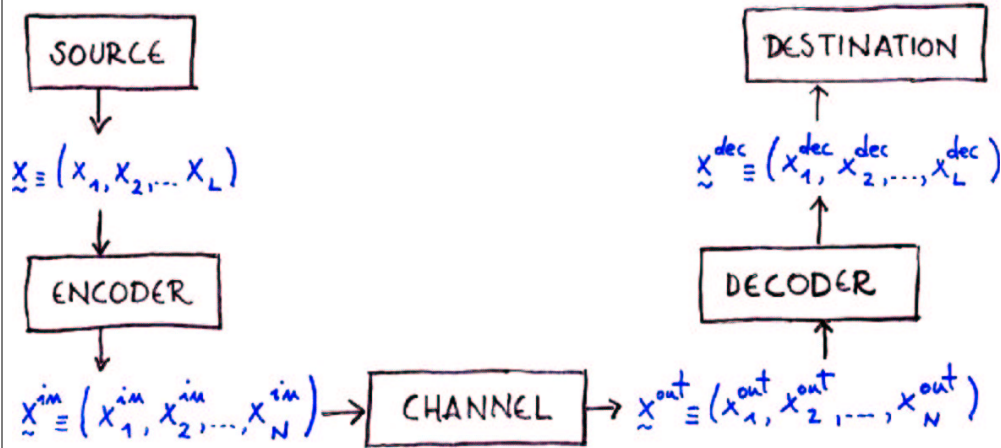
- example 2 : BINARY ERASURE CHANNEL (BEC)



$p \equiv$ "ERASURE PROBABILITY"

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ADD REDUNDANCY!

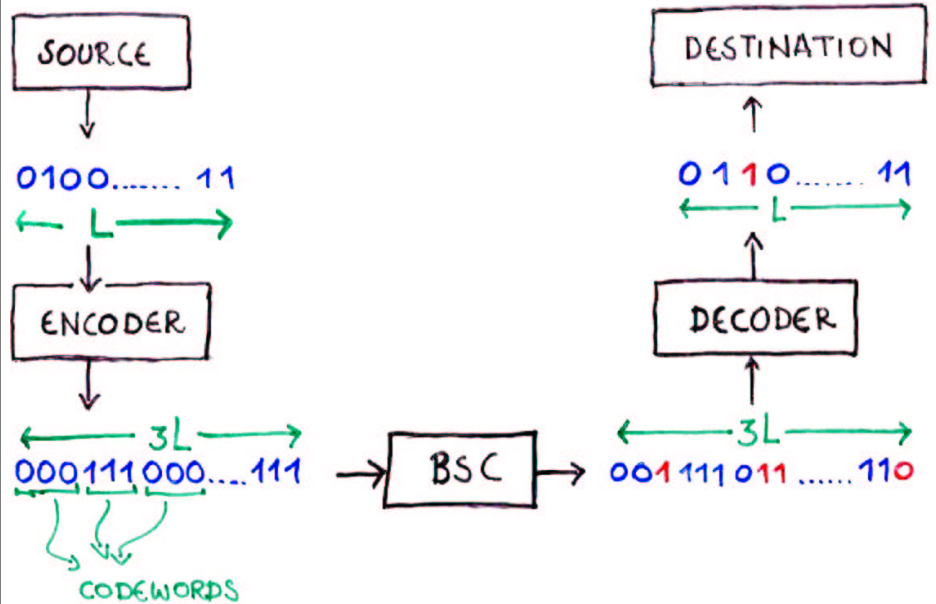


$$R \equiv \frac{L}{N} < 1$$

RATE

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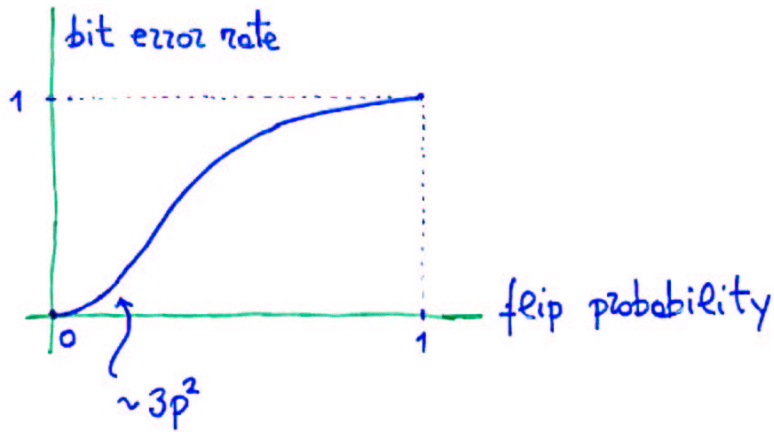
example : REPETITION CODE



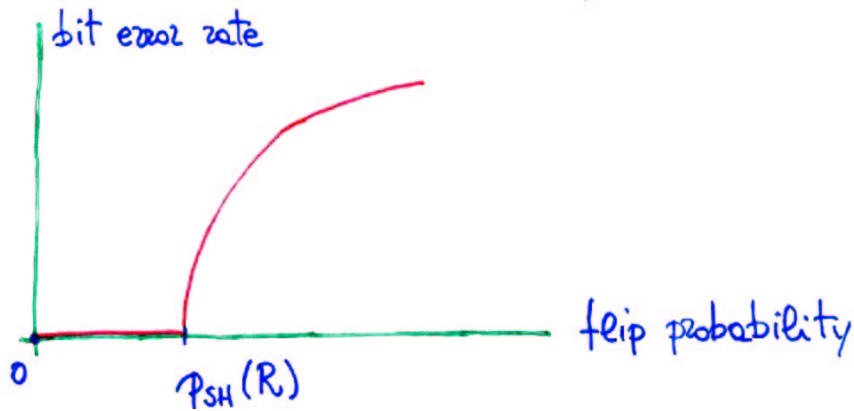
$$R = \frac{1}{3}$$

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PERFORMANCES



SHANNON SAID [1948]: "YOU CAN COMMUNICATE ERROR-FREE BELOW CAPACITY."



GOOD NEWS \rightarrow "ALMOST ANY CODE
ACHIEVE SHANNON BOUND"

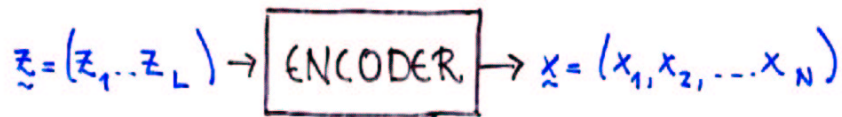
BAD NEWS \rightarrow "DECODING IS (IN GENERAL)
AN NP-COMplete PROBLEM"
[BERKELAMP,
MC ELIECEE,
VAN TILBOURG]

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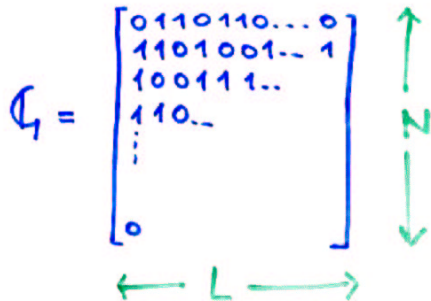
INTRODUCTION: LOW DENSITY PARITY CHECK CODES

[GALLAGER, MACKAY-NEAL, RICHARDSON-URBANKE,
LUBY, SPIELMAN, ...]

[KABASHIMA, SAAD, KANTER, A.M, SOURLAS...]



$$\underline{x} = G \underline{z} \pmod{2}$$

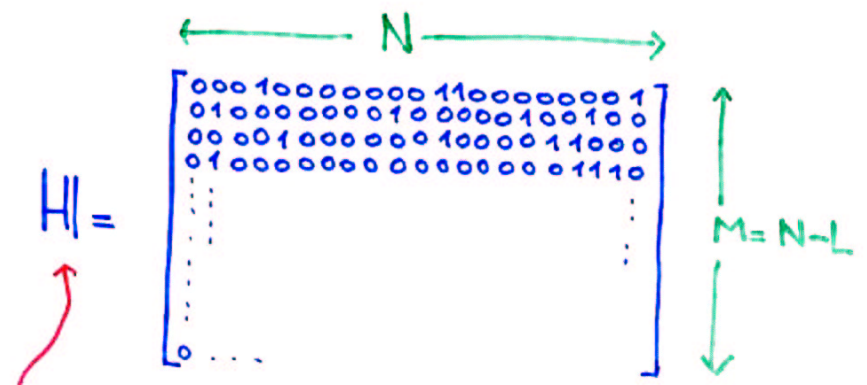


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CODEWORDS FORM A LINEAR SPACE :

$$\{ \underline{x} = G \underline{z} : \underline{z} \in \{0,1\}^L \} =$$

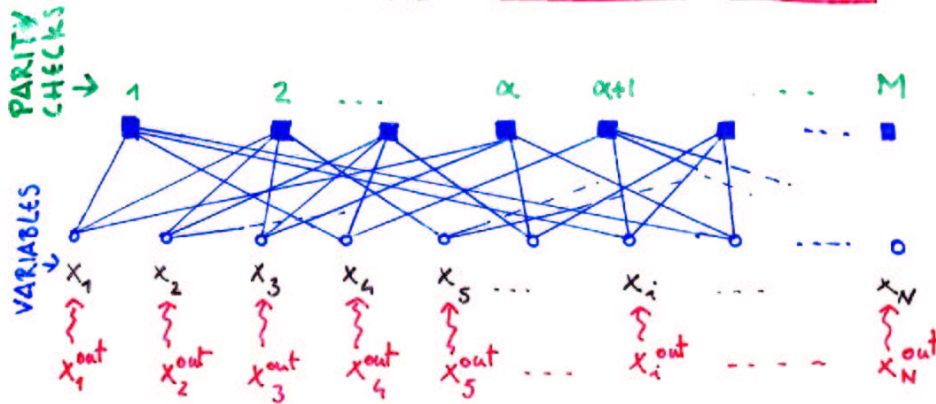
$$= \{ \underline{x} \in \{0,1\}^N : H \underline{x} = 0 \}$$



- PARITY CHECK MATRIX :
- RANDOM AND SPARSE
- k ONES PER ROW
 - e ONES PER COLUMN

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DECODING IS AN INFERENCE PROBLEM



$\alpha=2 \Rightarrow x_1 \oplus x_2 \oplus x_3 \oplus x_5 = 0$

$$P(\underline{x} | \underline{x}^{out}) = \frac{1}{Z} \prod_{\alpha=1}^M \psi(x_{\alpha_1}, \dots, x_{\alpha_k}) \prod_{i=1}^N p(x_i^{out} | x_i)$$

INTERACTIONS

RANDOM FIELD

PROBABILITY FOR \underline{x} TO BE THE CHANNEL INPUT

$$\psi(x_1, \dots, x_k) = \begin{cases} 1 & \text{if } x_1 \oplus x_2 \oplus \dots \oplus x_k = 0 \\ 0 & \text{ow.} \end{cases}$$

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DECODING PROBLEM : COMPUTE

$$P(x_i | \underline{x}^{out}) \equiv \sum_{\{x_j | j \neq i\}} P(\underline{x} | \underline{x}^{out})$$

EXPONENTIAL NUMBER OF SUMMANDS!!

$$x_i^{dec} = \arg \max_{x_i} P(x_i | \underline{x}^{out})$$

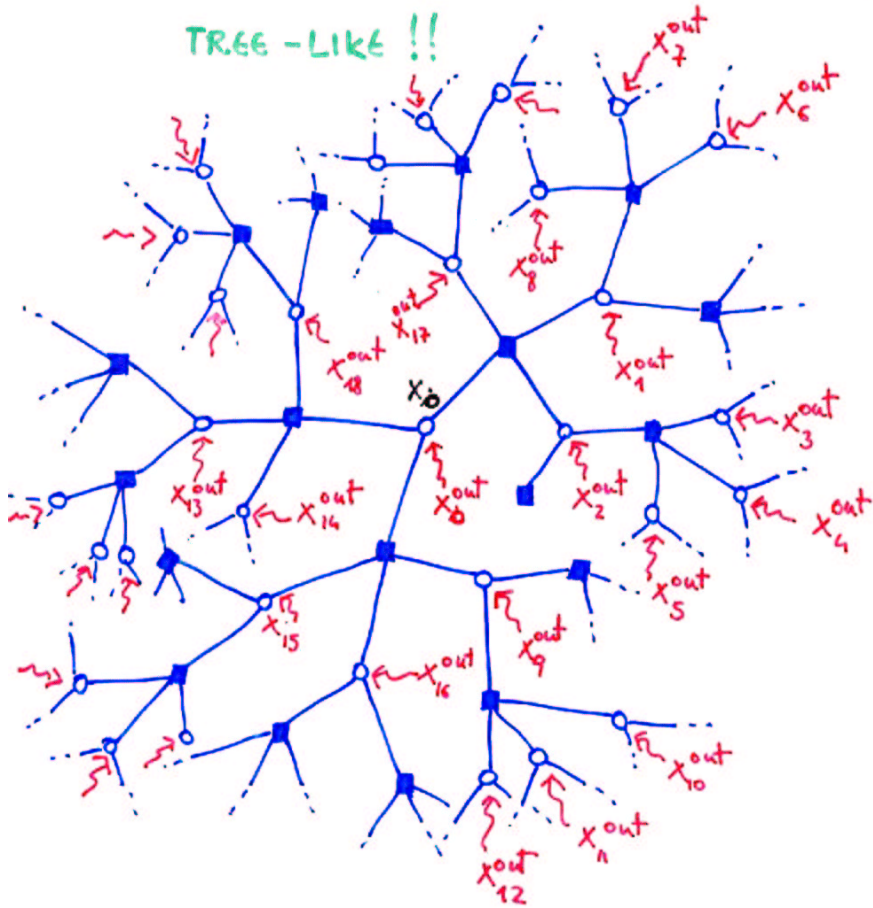
↑
DECODER
OUTPUT

[MAP DECODING]

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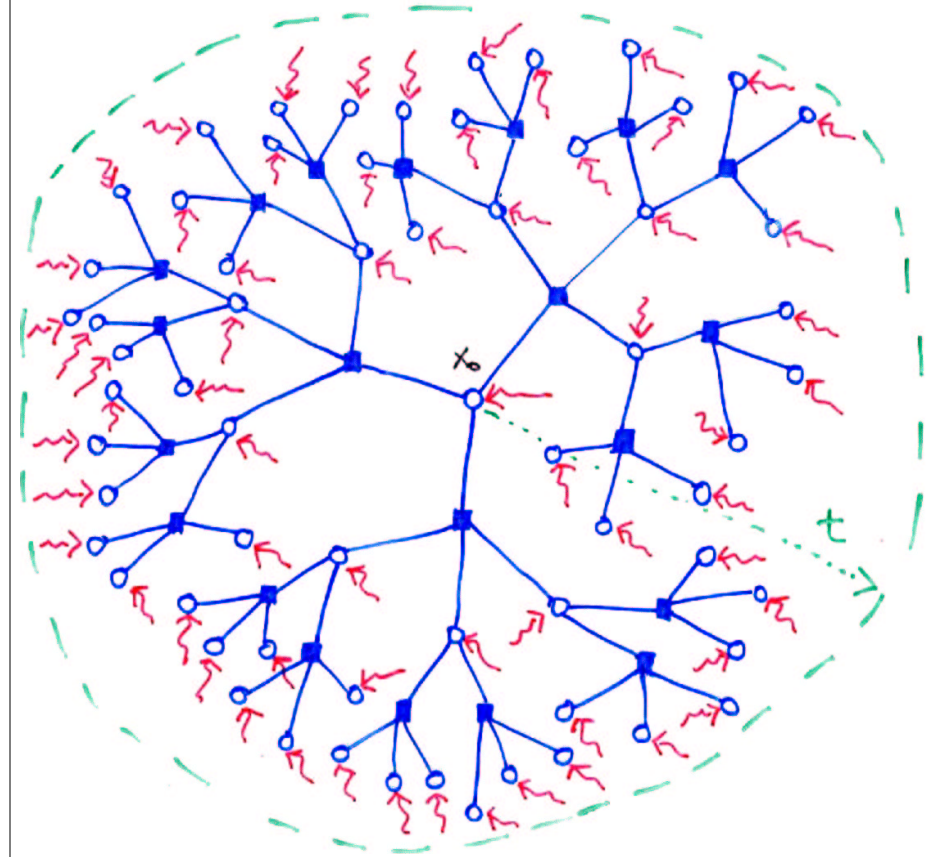
ITERATIVE DECODING

THE INTERACTION GRAPH IS LOCALLY TREE-LIKE !!



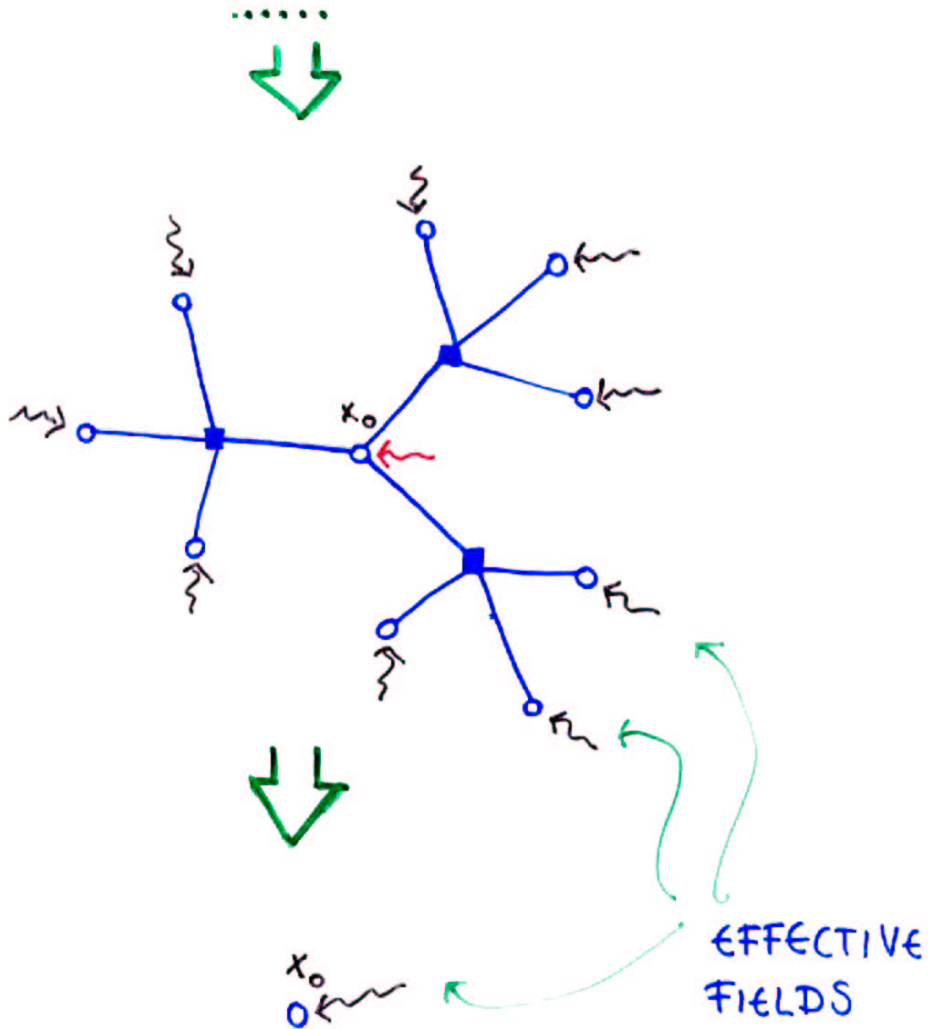
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1) CUT THE GRAPH AT A DISTANCE t FROM THE ROOT.....



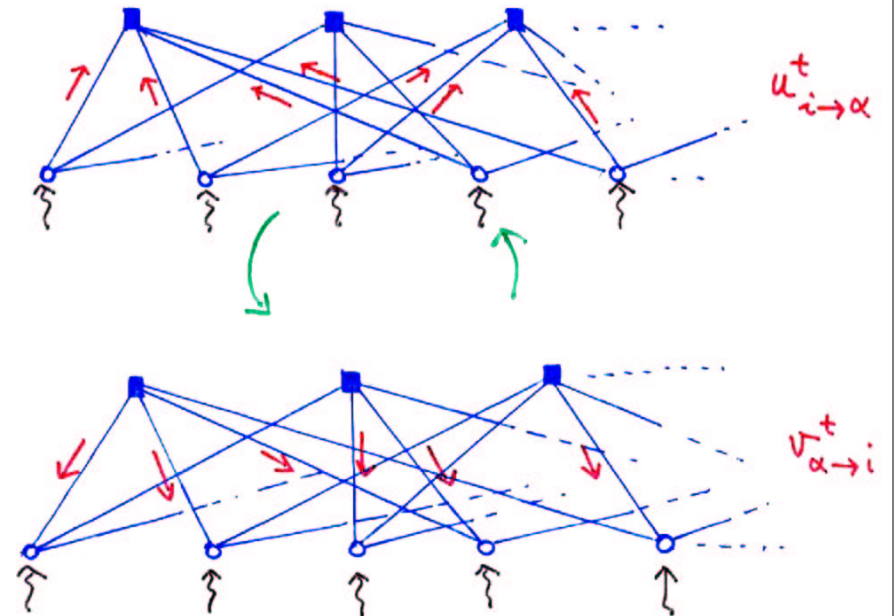
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2)...AND SUM RECURSIVELY OVER THE LEAVES



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MESSAGE-PASSING SCHEDULING
(BELIEF PROPAGATION)



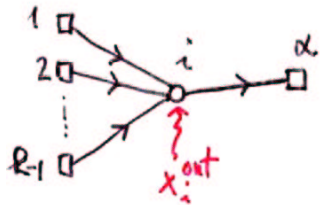
$$\{u_{i \rightarrow \alpha}^t, v_{\alpha \rightarrow i}^t\} \in \mathbb{R}$$

MESSAGES

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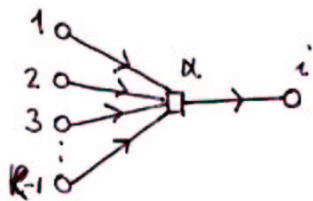
MESSAGE PASSING: UPDATES AT THE NODES

• VARIABLE NODES



$$u_{i \rightarrow \alpha}^{t+1} = h(x_i^{\text{out}}) + \sum_{\beta=1}^{R-1} v_{\beta \rightarrow i}^t$$

• CHECK NODES



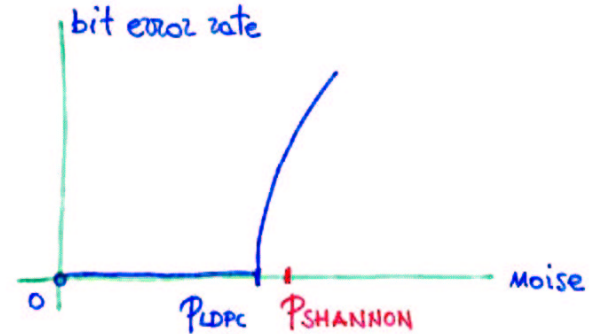
$$v_{\alpha \rightarrow i}^{t+1} = \text{arctanh} \left[\prod_{j=1}^{k-1} \tanh u_{j \rightarrow \alpha}^t \right]$$

CAVITY EQUATIONS

[≡ SADDLE POINT EQUATIONS OF THE BETHE FREE ENERGY $\phi(\underline{u}, \underline{v})$]

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PERFORMANCES



AMAZINGLY CLOSE TO SHANNON BOUND !!!

[cf. CHUNG, FORNEY, RICHARDSON URBANKE]

AND THEORETICAL QUESTIONS

- DOES THE ALGORITHM COMPUTE THE CORRECT MARGINALS ?
- WOULD MAP-DECODING IMPROVE THE PERFORMANCES ?
- COULD A DIFFERENT LOW-COMPLEXITY

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A LITTLE BIT OF PHYSICS

$$P(\underline{x} | \underline{x}^{out}) = \frac{1}{Z} \prod_{\alpha} \psi(x_{\alpha_1} \dots x_{\alpha_k}) \prod_i \phi(x_i^{out} | x_i)$$

$$H(\sigma) = -\beta \sum_{\alpha} \sigma_{\alpha_1} \dots \sigma_{\alpha_k} - \sum_{i=1}^N h_i \sigma_i$$

FERROMAGNETIC
MULTI-BODY
INTERACTIONS
($\beta \rightarrow \infty$)

RANDOM
MAGNETIC
FIELD

$$\sigma_i = (-1)^{x_i}$$

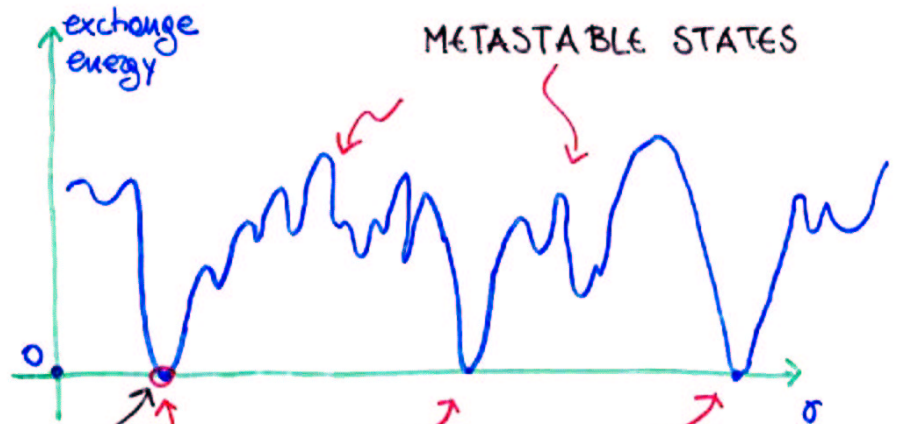
"SELF-INDUCED,"
FROSTRATION

LOW NOISE $\leftrightarrow h_i > 0$
HIGH NOISE $\leftrightarrow h_i < 0$

LESS BIASED | MORE BIASED

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THE HANDWAVING PICTURE:
WITHOUT EXTERNAL FIELD



2^{NR} CODEWORDS

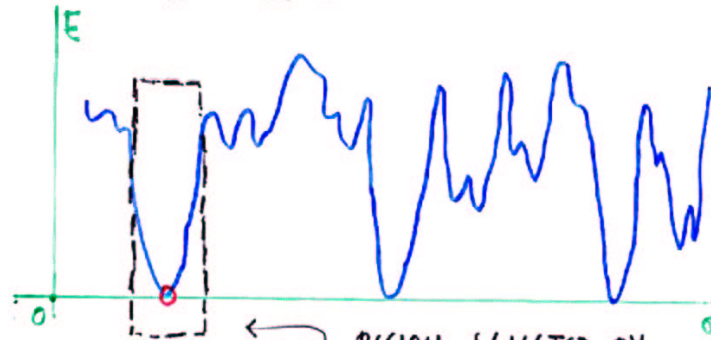
(\sim SOLUTIONS OF $H|\underline{x}=0$)

TRANSMITTED CODEWORD

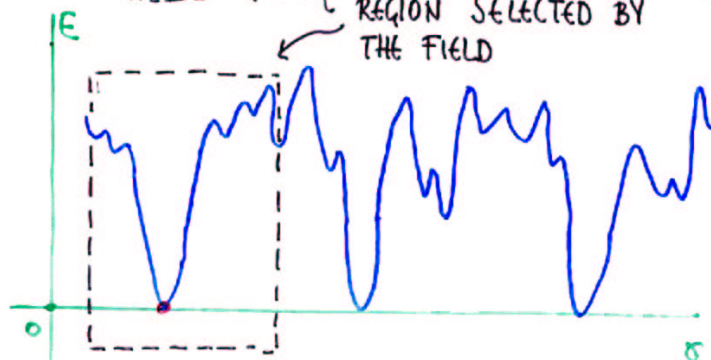
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THE HANDWAVING PICTURE :
WITH EXTERNAL FIELD

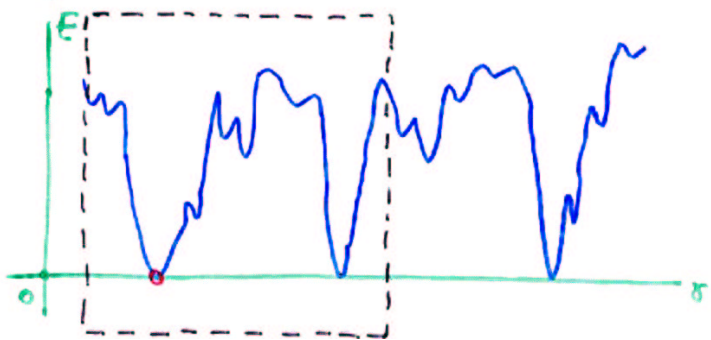
LOW-NOISE :
 $P > P_d$



INTERMEDIATE NOISE
 $P_d < P < P_c$

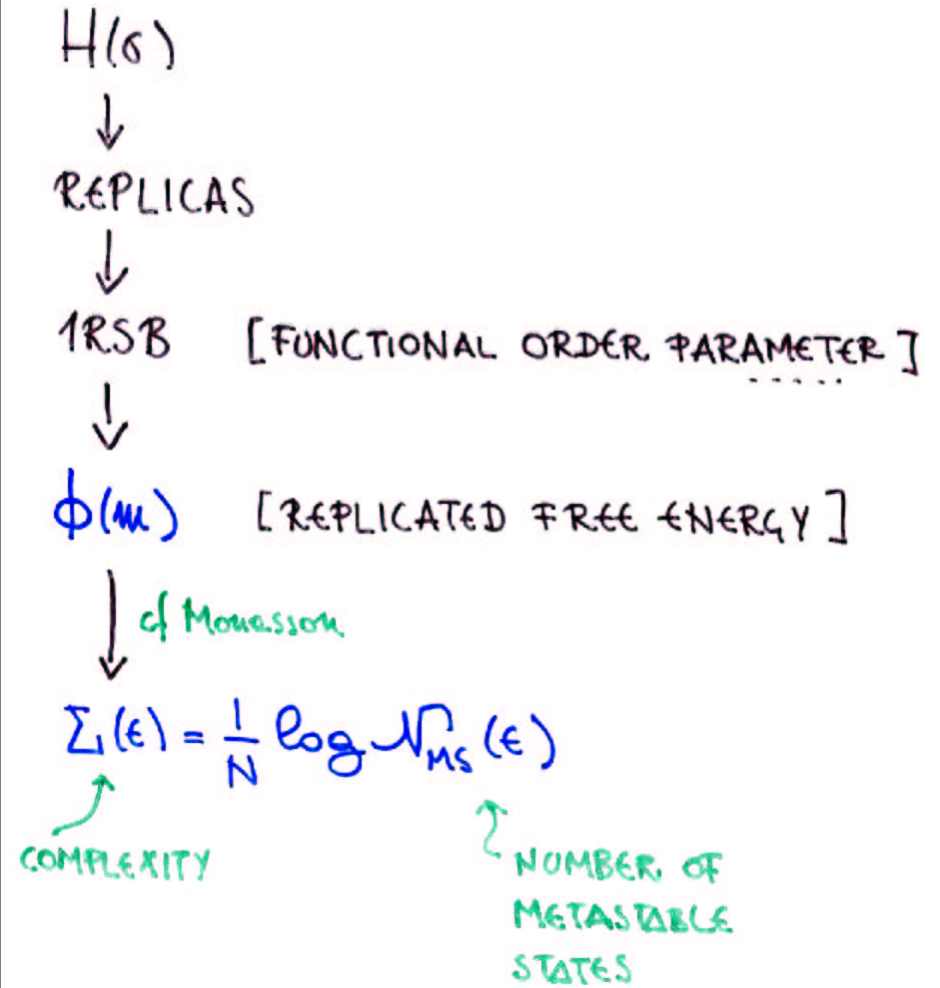


HIGH NOISE
 $P_c < P$



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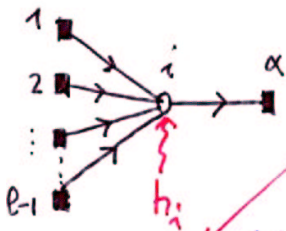
THE CALCULATION



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THE CALCULATION [AN ALTERNATIVE POINT OF VIEW]

LET US COUNT THE NUMBER OF FIXED POINTS OF THE ALGORITHM



NUMBER OF FIXED POINTS S.T. $u_{i \rightarrow \alpha} = u$ AND WITH FREE ENERGY F

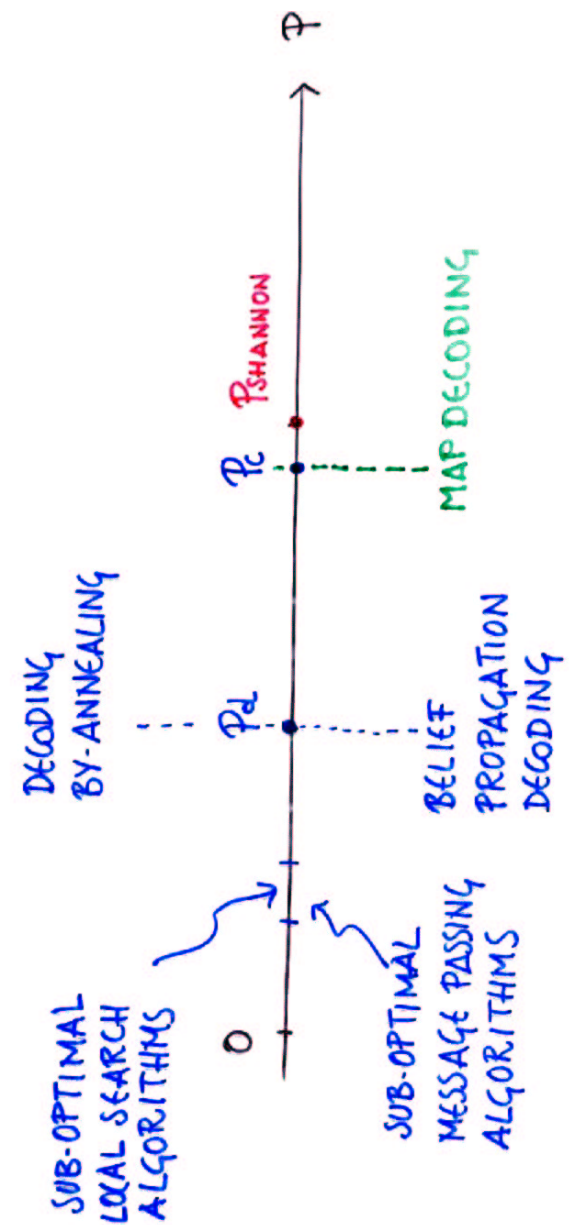
NUMBER OF ...

$$N_{i \rightarrow \alpha}(u; F) = \int \prod_{\beta=1}^{l-1} d\mathcal{W}_{\beta \rightarrow i}(v_{\beta}; F) \cdot \delta(u - h_i - \sum_{\beta} v_{\beta}) \cdot \delta(F - \sum_{\beta} F_{\beta} - \Delta F(\{v_{\beta}\}))$$

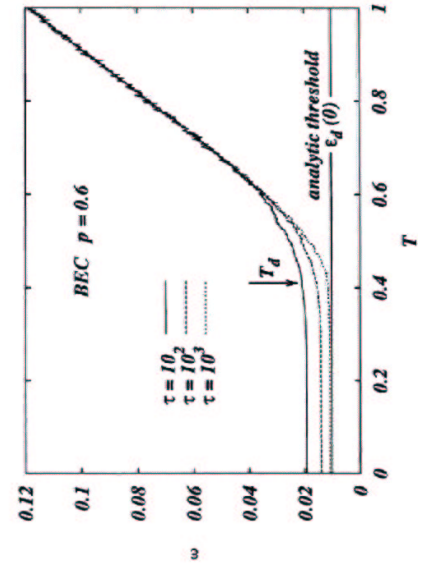
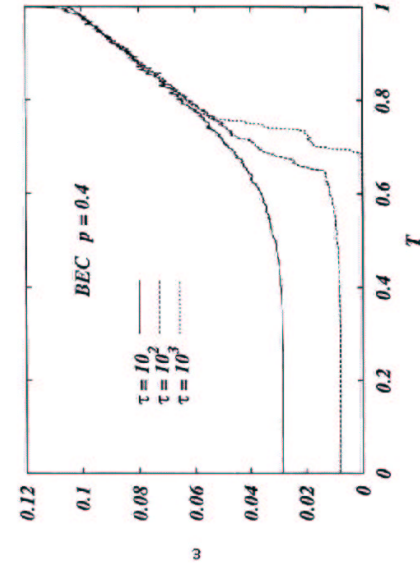
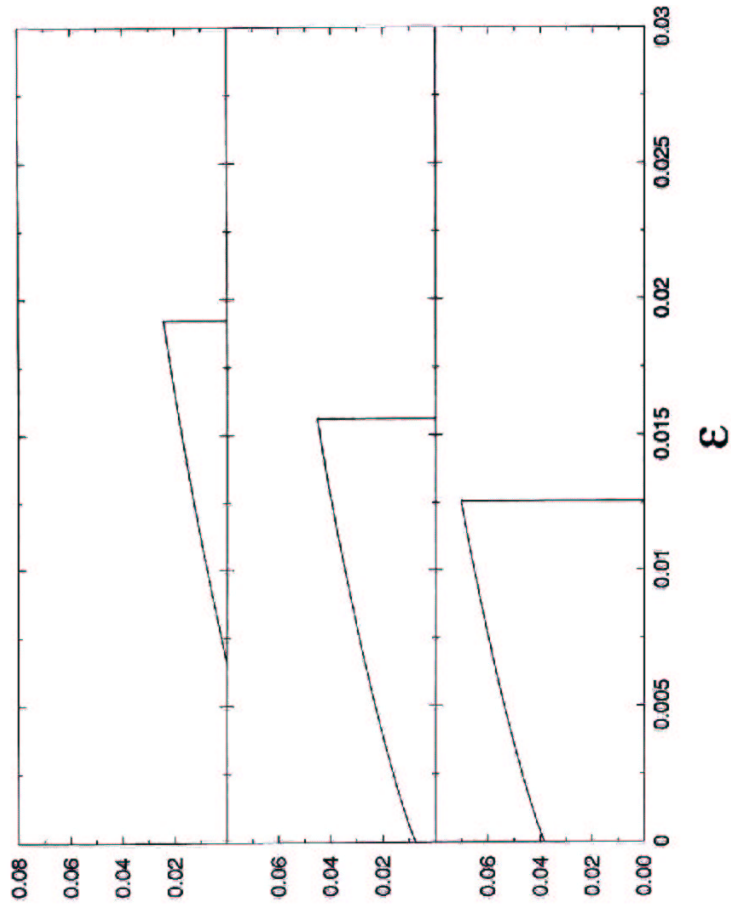
FIXED POINT CONDITION

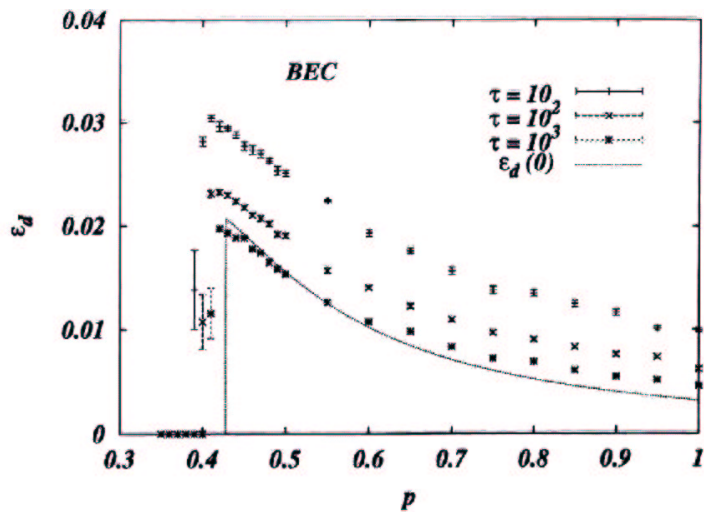
[PLUS A SIMILAR EQUATION FOR THE CHECK NODES UPDATING]

THE PHASE DIAGRAM



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DECODING - BY - ANNEALING

$$H(\sigma) = -\beta \sum_{\alpha} \sigma_{\alpha_1} \dots \sigma_{\alpha_k} - \sum_{i=1}^N h_i \sigma_i$$

- 1 • SET $\beta = 0$
- 2 • RUN A MONTE CARLO ALGORITHM FOR M STEPS
- 3 • IF $\beta < \beta_{\max}$, SET $\beta \leftarrow \beta + \Delta\beta$ AND GOTO 2.
- 4 • SET $\sigma_i^{\text{dec}} = \sigma_i$

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CONCLUSION

[OR: WHAT PHYSICS CAN LEARN FROM
THIS STORY]

- SINGLE -SAMPLE POINT OF VIEW
- THINK TO PHYSICAL APPROXIMATIONS
'ALGORITHMICALLY'