# Spectral Methods and Algorithms: applications in neuroscience 

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## Outline

1 Introduction and Motivation
2 Algebraic Graph Theory
3 Network Reorderings
■ Directed Hierarchies
4 Matrix Functions and Walks
■ Approximate Bipartite Substructures
■ Weighted Networks

- Strokes Vs Controls

5 Connecting it all Together

## Background

## Joel E. Cohen <br> Mathematics is biology's next microscope, only better

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Quantative/computational work in biology may be data driven or may arise through modelling

## Model

## Quantative, simplified description of a natural system

Useful for
e testing/comparing hypotheses
© making predictions
This talk will focus on networks: extracting useful information and modelling

## Typical Tasks

## Data Driven:

■ find well-connected clusters

- find specific connectivity substructures
- find 'important' nodes or links
- compare the properties of one network with another


## Modelling Arguments:

- summarize a network in terms of a few parameters
- explain how the connectivity has arisen
- discover missing or spurious links
- make predictions concerning future growth of the network


## Network Science: connections are important

Complex networks are the structural skeletons of complex systems


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LAlgebraic Graph Theory

## Graph Spectra

## Spectral methods:

© matrix representation of the network
e study the spectra of the resulting matrix, i.e., eigenvalues and eigenvectors

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Importantly:
e this allows us to compute graph invariants using basic linear algebra; and
e to implement data-mining tools to study networks, i.e., determine patterns and features

## Graph Spectra: adjacency matrix



$$
A= \begin{cases}1, & i \sim j \\ 0, & \text { otherwise }\end{cases}
$$

here $\sim$ denotes that vertices $i$ and $j$ are adjacent

LAlgebraic Graph Theory

## Other Possibilities Include ...



Graph Laplacian

$$
L= \begin{cases}d_{i}, & i=j \\ -a_{i j}, & \text { otherwise }\end{cases}
$$

here $d_{i}$ denotes the degree of node $i$

## Other Possibilities Include . . .

## Normalised Laplacian:

$$
\mathcal{L}= \begin{cases}1, & i=j \\ -\frac{a_{j i}}{\sqrt{d_{i j} d_{j}}}, & \text { otherwise }\end{cases}
$$

Signless Laplacian:

$$
Q= \begin{cases}d_{i}, & i=j \\ a_{i j}, & \text { otherwise }\end{cases}
$$

- less well studied
- can determine bipartite structures (Kirkland and Paul 2011)


## Line Graph:

■ used to detect community structure (Evans \& Lambiotte 2010)

## An Example: neuronal network of $C$. elegans

© C. elegans are tiny ( 1 mm long), transparent, round worms

■ Model organism in biology
■ Connectome consists of some 302 neurons linked by over 7000 synaptic connections

## Motivation



## Local C. elegans Neuronal Network



Adjacency Matrix
(a)

Reordered Adjacency Matrix

© Reordering the frontal neurons of C . elegans using eigenvectors of the signless Laplacian $Q$ reveals bipartite substructures

## C. elegans Example Ctd.



## Properties of the Different Graph Spectra

Important: cospectral graphs are not necessarily isomorphic

LAlgebraic Graph Theory

## Properties of the Different Graph Spectra

Important: cospectral graphs are not necessarily isomorphic

## Example:



$$
\chi_{G}=\lambda^{5}-4 \lambda^{3}
$$

## Properties of the Different Graph Spectra

Properties the spectrum (eigenvalues) can and cannot distinguish:

| Matrix | \# edges | bipartite | \# components | \# bipartite <br> components |
| :---: | :---: | :---: | :---: | :---: |
| $A$ | Yes | Yes | No | No |
| $L$ | Yes | No | Yes | No |
| $\mathcal{L}$ | No | Yes | Yes | Yes |
| $Q$ | Yes | No | No | Yes |

## Properties of the Different Graph Spectra

© Which graphs are determined by their spectrum? (Van Dam \& Haemers (2003))

■ For almost all graphs this is an open question
■ Numerical simulations suggest that 'almost all graphs are’

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© Which graphs are determined by their spectrum? (Van Dam \& Haemers (2003))

■ For almost all graphs this is an open question
■ Numerical simulations suggest that 'almost all graphs are’
© A combination of eigenvalues and eigenvectors does the trick, e.g., subgraph centrality


$$
C_{\mathcal{S}}(i)=\sum_{k=1}^{n} e^{\lambda_{k}} \mathbf{x}_{i}^{[k]^{2}}
$$

$$
C_{\mathcal{S}}=\left(\begin{array}{l}
1.6905 \\
1.6905 \\
3.7622 \\
1.6905 \\
1.6905
\end{array}\right) \quad \& \quad\left(\begin{array}{l}
2.3811 \\
2.3811 \\
1.0000 \\
2.3811 \\
2.3811
\end{array}\right)
$$

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## An Example: short-range structure

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- Solve

$$
\min \sum_{\{p \in \mathcal{P}\}}\left(p_{i}-p_{j}\right)^{2} a_{i j}
$$

$\mathcal{P}$ denotes the set of permutations of the integers $\{1, \ldots, n\}$

## An Example: short-range structure

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$$

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- An approximate solution is given by the first, non-zero eigen--vector of $L=D-A$

L Network Reorderings

- Directed Hierarchies


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$L_{\text {Network Reorderings }}$
L Directed Hierarchies

## Directed Hierarchies


$\left\llcorner_{\text {Network Reorderings }}\right.$

## - Directed Hierarchies

## Directed Hierarchies


e In general it is possible to find such an ordering iff we have a DAG
$L_{\text {Network Reorderings }}$
$\left\llcorner_{\text {Directed Hierarchies }}\right.$

## Out Minus In Degree

## One-Sum Optimisation Problem

$$
\min _{p \in \mathcal{P}} \sum_{i, j}\left(p_{i}-p_{j}\right) a_{i j}
$$


$L_{\text {Network Reorderings }}$
L Directed Hierarchies

## Out Minus In Degree

## One-Sum Optimisation Problem

$$
\min _{p \in \mathcal{P}} \sum_{i, j}\left(p_{i}-p_{j}\right) a_{i j}
$$

## Proof

$$
\begin{aligned}
\sum_{i, j}\left(p_{j}-p_{i}\right) a_{i j} & =\sum_{i} p_{i} \cdot \operatorname{deg}_{i}^{\text {out }}-\sum_{j} p_{j} \cdot \operatorname{deg}_{j}^{\text {in }} \\
& =\sum_{i} p_{i} \cdot\left(\operatorname{deg}_{i}^{\text {out }}-\operatorname{deg}_{i}^{\text {in }}\right)
\end{aligned}
$$

$L_{\text {Network Reorderings }}$

- Directed Hierarchies


## Synthetic Network Example



e Hierarchical structure is uncovered in RHS using out-in degree

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$\left\llcorner_{\text {Matrix Functions and Walks }}\right.$

## Paths Vs Walks



Path


Walk

## Shortest Path

$L_{\text {Matrix Functions and Walks }}$

## Counting Walks

From the following identity

$$
\left(A^{k}\right)_{i j}=\sum_{i_{1}=1}^{n} \sum_{i_{2}=1}^{n} \cdots \sum_{i_{k-1}=1}^{n} a_{i, i_{1}} a_{i_{1}, i_{2}} \cdots a_{i_{k-1}, j}
$$

we see that $\left(A^{k}\right)_{i j}$ counts the number of different walks of length $k$ between nodes $i$ and $j$;
$\left\llcorner_{\text {Matrix Functions and Walks }}\right.$

## Counting Walks

## From the following identity


we see that $\left(A^{k}\right)_{i j}$ counts the number of different walks of length $k$
between nodes $i$ and $j$; moreover, the quantity

$$
F_{i j}=\left(c_{0} I+c_{1} A+c_{2} A^{2}+c_{3} A^{3} \cdots\right)_{i j}
$$

with the $c_{k}$ constant, gives a measure of the total number of walks between nodes $i$ and $j$

## Counting Walks: matrix functions

© For suitable choices of $\left\{c_{k}\right\}_{k \geq 0}$ the series overpage converges
■ for example, $c_{k}=\left\{\delta^{k}\right\}_{k \geq 0}$, gives the matrix resolvent $(\mathbf{I}-\delta \mathbf{A})^{-1}$
■ different choices of $c_{k}$ allow for different scalings

## Counting Walks: matrix functions

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■ for example, $\boldsymbol{c}_{k}=\left\{\delta^{k}\right\}_{k \geq 0}$, gives the matrix resolvent $(\mathbf{I}-\delta \mathbf{A})^{-\mathbf{1}}$

- different choices of $c_{k}$ allow for different scalings
© In particular, if $f$ is defined on the spectrum of $A$, we can define

$$
F=f(A)=P f(D) P^{-1}
$$

(here $\mathrm{A}=\mathrm{PDP}^{-1}$ is the eigendecomposition)
Proof Let

$$
f(A)=c_{0} I+\sum_{k=1}^{\infty} c_{k} A^{k}
$$

substituting $A=P D P^{-1}$ into the above

$$
f(A)=f\left(P D P^{-1}\right)=c_{0} I+\sum_{k=1}^{\infty} c_{k}\left(P D P^{-1}\right)^{k}=P\left(c_{0} I+\sum_{k=1}^{\infty} c_{k} D^{k}\right) P^{-1}
$$

$L_{\text {Matrix Functions and Walks }}$

## Example: communicability



$$
A=\left(\begin{array}{llll}
0 & 1 & 1 & 1 \\
1 & 0 & 1 & 0 \\
1 & 1 & 0 & 1 \\
1 & 0 & 1 & 0
\end{array}\right), A^{2}=\left(\begin{array}{llll}
3 & 1 & 2 & 1 \\
1 & 2 & 1 & 2 \\
2 & 1 & 3 & 1 \\
1 & 2 & 1 & 2
\end{array}\right), \cdots
$$

Communicability between distinct nodes $i$ and $j$

$$
\left(1+A+A^{2} / 2!+A^{3} / 3!+\cdots\right)_{i j}
$$

that is $\left(e^{A}\right)_{i j}$
Spectral form: $\left(\sum_{k=1}^{n} e^{\lambda_{k}} \mathbf{x}^{[k]}(i) \mathbf{x}^{[k]}(j)\right)_{i j}$

## Example: communicability


$\exp (\mathrm{A})$


Communicability applied to the frontal network of C. elegans
$\square_{\text {Matrix Functions and Walks }}$
$\square_{\text {Approximate Bipartite Substructures }}$

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## $\square_{\text {Approximate Bipartite Substructures }}$

## Approximate Directed Bipartite Community

Crofts, Estrada, Higham, Taylor Elec. Trans. Numer. Anal (2010)

Distinct subsets of nodes $S_{1}$ and $S_{2}$ such that

- $S_{1}$ has few internal links
- $S_{2}$ has few internal links
- there are many $S_{1} \rightarrow S_{2}$ links
- few other links involve $S_{1}$ or $S_{2}$

$L_{\text {Matrix Functions and Walks }}$
$\square_{\text {Approximate Bipartite Substructures }}$


## An alternating walk of length $k$ from node $i_{1}$ to node $i_{k+1}$

is a list of nodes

$$
i_{1}, i_{2}, i_{3}, \ldots, i_{k+1}
$$

such that $a_{i_{s}, i_{s+1}} \neq 0$ for $s$ odd, and $a_{i_{s+1}}, i_{s} \neq 0$ for $s$ even
Loosely, an alternating walk is a traversal that successively follows links in the forward and reverse directions

$\square_{\text {Matrix Functions and Walks }}$
$\square_{\text {Approximate Bipartite Substructures }}$

## This Motivates . . .

$$
f(A)=I-A+\frac{A A^{T}}{2!}-\frac{A A^{T} A}{3!}+\frac{A A^{T} A A^{T}}{4!}-\cdots
$$

Overall idea: $f(A)+f\left(A^{T}\right)$ has
■ positive values representing inter-community $S_{1} \leftrightarrow S_{1}$ and $S_{2} \leftrightarrow S_{2}$ relationships, and
■ negative values representing extra-community $S_{1} \leftrightarrow S_{2}$ relationships

Also, $f(A)+f\left(A^{T}\right)$ is a symmetric matrix, so amenable to standard clustering techniques

Note: $f(A)$ defined above is not a matrix function

UCSB 2011

- Matrix Functions and Walks
- Approximate Bipartite Substructures


## Synthetic Example



$$
n z=252
$$



Matrix Functions and Walks

- Approximate Bipartite Substructures


## C. elegans Neural Data

Automates the computations of Durbin, (PhD thesis, Cambridge, 1987)


Subnetwork of A


$\square_{\text {Matrix Functions and Walks }}$
-Weighted Networks

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$L_{\text {Matrix Functions and Walks }}$
-Weighted Networks

## Communicability for a Weighted Network

W is symmetric with non-negative real weights
Let

$$
d_{i}=\sum_{k=1}^{n} w_{i k} \quad \& \quad D:=\operatorname{diag}\left(d_{i}\right)
$$

Normalisation:

$$
W \mapsto D^{-1 / 2} W D^{-1 / 2}
$$

for at least 2 reasons
■ avoid overflow

- to stifle promiscuous nodes
$L_{\text {Matrix Functions and Walks }}$
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Communicability measure:

$$
\exp \left(D^{-1 / 2} W D^{-1 / 2}\right)
$$

$L_{\text {Matrix Functions and Walks }}$

## - Weighted Networks

## But What Does it Actually Mean?



$$
W=\left(\begin{array}{cccc}
0 & w_{12} & w_{13} & w_{14} \\
w_{21} & 0 & w_{23} & w_{24} \\
w_{31} & w_{32} & 0 & w_{34} \\
w_{41} & w_{42} & w_{43} & 0
\end{array}\right)
$$

The $k$ th powers of $W$ provide a measure of the total strength contained within walks of length $k$ between nodes $i$ and $j$ :

$$
\left(W^{2}\right)_{i j}=\sum_{k} w_{i k} w_{k j},\left(W^{3}\right)_{i j}=\cdots
$$

$\square_{\text {Matrix Functions and Walks }}$
LWeighted Networks

## Example: anatomical connectivity data

© 9 subjects - at least 6 months following first, left-hemisphere, subcortical stroke; and 10 (18) age matched controls
© Diffusion Tensor Imaging computes all connections between all voxels
© Connectivity network based on the Harvard-Oxford cortical and subcortical structural atlas: 48 cortical regions and 8 subcortical regions

http://wwww.fmrib.ox.ac.uk/fsl/fslview/
atlas-descriptions.html

- Matrix Functions and Walks


## - Weighted Networks

## Unsupervised Clustering of Patients




Crofts \& Higham, Roy. Soc. Interface (2009)
$\square_{\text {Matrix Functions and Walks }}$
-Weighted Networks

## Communicability Adds Value to the Raw Data



Communicability (Control)

$L_{\text {Matrix Functions and Walks }}$

- Weighted Networks


## What Are we Actually Detecting?

## Question

Exactly what enables us to differentiate between strokes \& controls?

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## Question

Exactly what enables us to differentiate between strokes \& controls?
© If it is merely the fact that many connections have been destroyed close to the infarcted region this is not very interesting - an MRI scan can tell us this with no further analysis needed!

- Why not rerun the tractographies including only those connections within the non-stroke hemisphere?
■ If we still distinguish between the two classes, this will be more relevant from a bio point of view
- Matrix Functions and Walks

LWeighted Networks

## RHS Hemisphere Sorted by Brain Region





$\square_{\text {Matrix Functions and Walks }}$

## -Weighted Networks

## Left (Stroke Side) Vs Right


$\square_{\text {Matrix Functions and Walks }}$
-Weighted Networks

## Overlay of Stroke Overlap Volume



Highlighted regions are those where patients showed reduced communicability relative to controls

## UCSB 2011

Matrix Functions and Walks
-Weighted Networks

## Crofts et al. Neurolmage (2011)



And increased communicability
$L_{\text {Matrix Functions and Walks }}$

- Weighted Networks


## GSVD

© Generalised singular value decomposition

$$
A=U C X^{-1} \quad \text { and } \quad B=V S X^{-1}
$$

of a pair of matrices can be used to determine clusters that are 'good' in one network and 'poor' in the other - and vice versa [Xiao et al. (2011)]
© Another way to understand this is to note that, in the case where $A$ and $B$ are invertible the GSVD is closely related to the SVD of $A B^{-1}$ and $B A^{-1}$, hence

$$
A B^{-1}=U C S^{-1} V^{T} \quad \text { and } \quad B A^{-1}=V S C^{-1} U^{T}
$$

## GSVD Example: anatomical connectivity data

Controls


Stroke


GSVD finds a group of brain regions that are much better connected in controls than strokes

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## The Take Home Message

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- scalable
- most work has focused on the adjacency matrix


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© Walk based measures are attractive since:
■ combinatorics are described in terms of basic linear algebra
■ information does not necessarily flow along geodesics
- walks are more tolerent to errors


## The Take Home Message

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- scalable
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© Walk based measures are attractive since:
■ combinatorics are described in terms of basic linear algebra
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e GSVD and extensions:
- tensor decompositions (plus multiple tasks/modalities, time-dependent networks)
■ NNMF (Lee et al (2010)), ICA (Smith et al (2005))


## Thank you!

## Colleagues at Strathclyde: Des Higham, Ernesto Estrada \& Alan Taylor

Colleagues from Oxford: Heidi Johanson-Berg \& Tim Behrens The stroke data was supplied by Rose Bosnell

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