

-Welcome

Spectral Methods and Algorithms: applications in neuroscience

Jonathan J. Crofts

Department of Physics & Mathematics Nottingham Trent University

August 5, 2011

jonathan.crofts@ntu.ac.uk

J. J. Crofts ()

Outline

Introduction and Motivation
 Algebraic Graph Theory
 Network Reorderings

 Directed Hierarchies

 Matrix Functions and Walks

 Approximate Bipartite Substructures
 Weighted Networks

 Strokes Vs Controls

 Connecting it all Together

-Introduction and Motivation

Background

Joel E. Cohen

Mathematics is biology's next microscope, only better

-Introduction and Motivation

Background

Joel E. Cohen

Mathematics is biology's next microscope, only better

Quantative/computational work in biology may be **data driven** or may arise through **modelling**

Model

Quantative, simplified description of a natural system

Useful for

- e testing/comparing hypotheses
- e making predictions

This talk will focus on **networks**: **extracting useful information** and **modelling**

J. J. Crofts ()

Typical Tasks

Data Driven:

- find well-connected clusters
- find specific connectivity substructures
- find 'important' nodes or links
- **compare** the properties of one network with another

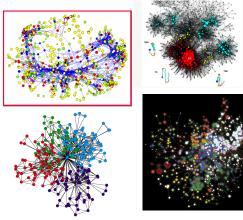
Modelling Arguments:

- summarize a network in terms of a few parameters
- explain how the connectivity has arisen
- discover missing or spurious links
- make predictions concerning future growth of the network

-Introduction and Motivation

Network Science: connections are important

Complex networks are the structural skeletons of complex systems



Outline

- 1 Introduction and Motivation
- 2 Algebraic Graph Theory
- **3** Network Reorderings
 - Directed Hierarchies
- 4 Matrix Functions and Walks
 - Approximate Bipartite Substructures
 - Weighted Networks
 - Strokes Vs Controls
- 5 Connecting it all Together

Graph Spectra

Spectral methods:

- matrix representation of the network
- study the spectra of the resulting matrix, i.e., eigenvalues and eigenvectors



Graph Spectra

Spectral methods:

matrix representation of the network

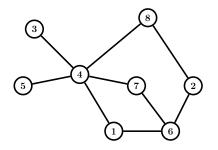


study the spectra of the resulting matrix, i.e., eigenvalues and eigenvectors

Importantly:

- this allows us to compute graph invariants using basic linear algebra; and
- to implement data-mining tools to study networks, i.e., determine patterns and features

Graph Spectra: adjacency matrix



	/ 0	0	0	1	0	1	0	0	
	0	0	0	0	0	1	0	1	Ì
	0	0	0	1	0	0	0	0	
Λ	1	0	1	0	1	0	1	1	
A =	0	0	0	1	0	0	0	0	
	1	1	0	0	0	0	1	0	
	0	0	0	1	0	1	0	0	
	\ 0	1	0	1	0	0	0 0 1 0 1 0 0	0)

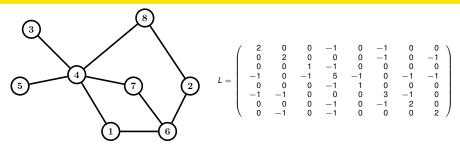
$$A = \begin{cases} 1, & i \sim j \\ 0, & \text{otherwise} \end{cases}$$

here ~ denotes that vertices *i* and *j* are **adjacent**

J. J. Crofts ()

Algebraic Graph Theory

Other Possibilities Include ...



Graph Laplacian

$$L = \begin{cases} d_i, & i = j \\ -a_{ij}, & \text{otherwise} \end{cases}$$

here d_i denotes the **degree** of node *i*

J. J. Crofts ()

- Algebraic Graph Theory

Other Possibilities Include ...

Normalised Laplacian:

$$\mathcal{L} = egin{cases} 1, & i = j \ -rac{a_{ij}}{\sqrt{d_i d_j}}, & ext{otherwise} \end{cases}$$

Signless Laplacian:

$$Q = \begin{cases} d_i, & i = j \\ a_{ij}, & \text{otherwise} \end{cases}$$

less well studied

can determine bipartite structures (Kirkland and Paul 2011)

Line Graph:

used to detect community structure (Evans & Lambiotte 2010)

J. J. Crofts ()

UCSB 2011

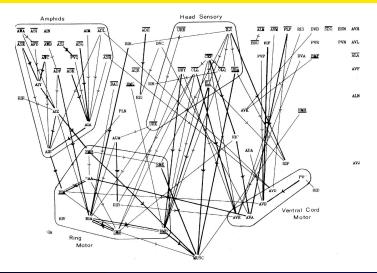
An Example: neuronal network of C. elegans

- C. elegans are tiny (1mm long), transparent, round worms
 - Model organism in biology
 - Connectome consists of some 302 neurons linked by over 7000 synaptic connections



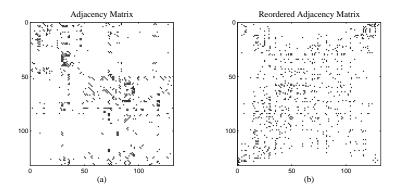
Algebraic Graph Theory

Motivation



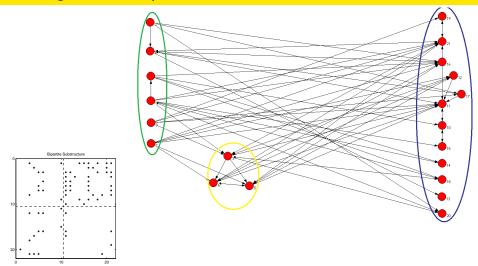
J. J. Crofts ()

Local C. elegans Neuronal Network



Reordering the frontal neurons of C. elegans using eigenvectors of the signless Laplacian Q reveals bipartite substructures

C. elegans Example Ctd.

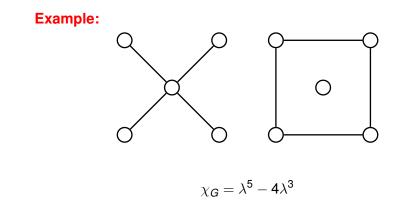


Properties of the Different Graph Spectra

Important: cospectral graphs are not necessarily isomorphic

Properties of the Different Graph Spectra

Important: cospectral graphs are not necessarily isomorphic



Properties of the Different Graph Spectra

Properties the spectrum (eigenvalues) can and cannot distinguish:

Matrix	# edges	bipartite	# components	# bipartite	
				components	
A	Yes	Yes	No	No	
L	Yes	No	Yes	No	
\mathcal{L}	No	Yes	Yes	Yes	
Q	Yes	No	No	Yes	

- Algebraic Graph Theory

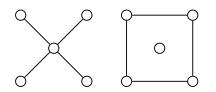
Properties of the Different Graph Spectra

- Which graphs are determined by their spectrum? (Van Dam & Haemers (2003))
 - For almost all graphs this is an open question
 - Numerical simulations suggest that 'almost all graphs are'

- Algebraic Graph Theory

Properties of the Different Graph Spectra

- Which graphs are determined by their spectrum? (Van Dam & Haemers (2003))
 - For almost all graphs this is an open question
 - Numerical simulations suggest that 'almost all graphs are'
- A combination of eigenvalues and eigenvectors does the trick, e.g., subgraph centrality



$$C_{S}(i) = \sum_{k=1}^{n} e^{\lambda_{k}} \mathbf{x}_{i}^{[k]^{2}}$$
$$= \begin{pmatrix} 1.6905\\ 1.6905\\ 3.7622\\ 1.6905\\ 1.6905 \end{pmatrix} \& \begin{pmatrix} 2.3811\\ 2.3811\\ 2.3811\\ 2.3811\\ 2.3811 \end{pmatrix}$$

 $C_{S} =$

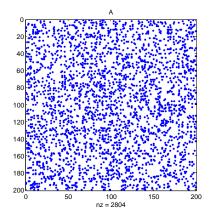
Outline

- 1 Introduction and Motivation
- 3 Network Reorderings
 - Directed Hierarchies
- 4 Matrix Functions and Walks
 - Approximate Bipartite Substructures
 - Weighted Networks
 - Strokes Vs Controls
- 5 Connecting it all Together

-Network Reorderings

An Example: short-range structure

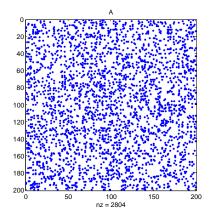
[©] The network reordering problem:



- Network Reorderings

An Example: short-range structure

[©] The network reordering problem:



Solve

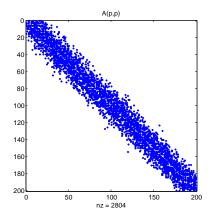
$$\min\sum_{\{\boldsymbol{p}\in\mathcal{P}\}}(\boldsymbol{p}_i-\boldsymbol{p}_j)^2\boldsymbol{a}_{ij}$$

 \mathcal{P} denotes the set of permutations of the integers $\{1, \ldots, n\}$

- Network Reorderings

An Example: short-range structure

[©] The network reordering problem:



Solve

$$\min\sum_{\{\boldsymbol{p}\in\mathcal{P}\}}(\boldsymbol{p}_i-\boldsymbol{p}_j)^2\boldsymbol{a}_{ij}$$

 \mathcal{P} denotes the set of permutations of the integers $\{1, \ldots, n\}$

• An **approximate solution** is given by the first, non-zero eigen--vector of L = D - A

-Network Reorderings

Directed Hierarchies

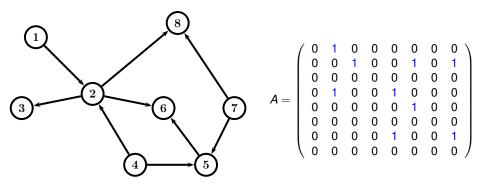
Outline

- 1 Introduction and Motivation
- 2 Algebraic Graph Theory
- 3 Network Reorderings
 - Directed Hierarchies
- 4 Matrix Functions and Walks
 - Approximate Bipartite Substructures
 - Weighted Networks
 - Strokes Vs Controls
- 5 Connecting it all Together

- Network Reorderings

Directed Hierarchies

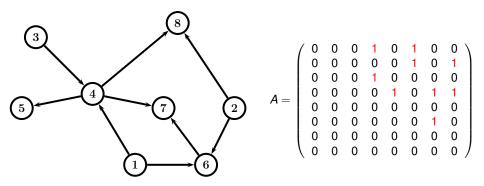
Directed Hierarchies



- Network Reorderings

Directed Hierarchies

Directed Hierarchies



In general it is possible to find such an ordering iff we have a DAG

-Network Reorderings

Directed Hierarchies

Out Minus In Degree

One-Sum Optimisation Problem

$$\min_{oldsymbol{
ho}\in\mathcal{P}}\sum_{i,j}(oldsymbol{
ho}_i-oldsymbol{
ho}_j)oldsymbol{a}_{ij}$$

Proof

$$\sum_{i,j} (p_j - p_i) a_{ij} = \sum_i p_i \cdot \deg_i^{\text{out}} - \sum_j p_j \cdot \deg_j^{\text{in}}$$
$$= \sum_i p_i \cdot (\deg_i^{\text{out}} - \deg_i^{\text{in}})$$

-Network Reorderings

Directed Hierarchies

Out Minus In Degree

One-Sum Optimisation Problem

$$\min_{oldsymbol{
ho}\in\mathcal{P}}\sum_{i,j}(oldsymbol{
ho}_i-oldsymbol{
ho}_j)oldsymbol{a}_{ij}$$

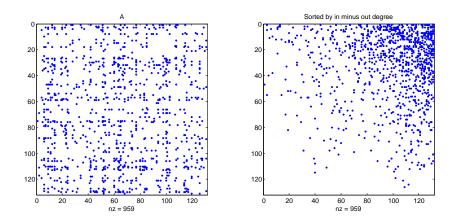
Proof

$$\sum_{i,j} (p_j - p_i) a_{ij} = \sum_i p_i \cdot \deg_i^{\text{out}} - \sum_j p_j \cdot \deg_j^{\text{in}}$$
$$= \sum_i p_i \cdot (\deg_i^{\text{out}} - \deg_i^{\text{in}})$$

-Network Reorderings

- Directed Hierarchies

Synthetic Network Example



[®] Hierarchical structure is uncovered in RHS using out-in degree

J. J. Crofts ()

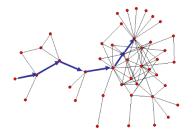
Outline

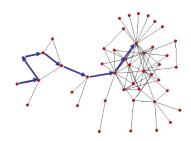
- 1 Introduction and Motivation
- 2 Algebraic Graph Theory
- 3 Network Reorderings
 - Directed Hierarchies
- 4 Matrix Functions and Walks
 - Approximate Bipartite Substructures
 - Weighted Networks
 - Strokes Vs Controls

5 Connecting it all Together

Matrix Functions and Walks

Paths Vs Walks





Path

Walk

Shortest Path

J. J. Crofts ()

UCSB 2011

-Matrix Functions and Walks

Counting Walks

From the following identity

$$(\mathbf{A}^{k})_{ij} = \sum_{i_{1}=1}^{n} \sum_{i_{2}=1}^{n} \cdots \sum_{i_{k-1}=1}^{n} a_{i,i_{1}} a_{i_{1},i_{2}} \cdots a_{i_{k-1},j},$$

we see that $(A^k)_{ij}$ counts the number of different walks of length k between nodes *i* and *j*;

$$F_{ij} = (c_0 I + c_1 A + c_2 A^2 + c_3 A^3 \cdots)_{ij}$$

with the *c_k* constant, gives a **measure of the total number of walks** between nodes *i* and *j* -Matrix Functions and Walks

Counting Walks

From the following identity

$$(A^k)_{ij} = \sum_{i_1=1}^n \sum_{i_2=1}^n \cdots \sum_{i_{k-1}=1}^n a_{i_1,i_1} a_{i_1,i_2} \cdots a_{i_{k-1},j_k}$$

we see that $(A^k)_{ij}$ counts the number of different walks of length k between nodes *i* and *j*; moreover, the quantity

$$F_{ij} = (c_0 I + c_1 A + c_2 A^2 + c_3 A^3 \cdots)_{ij}$$

with the c_k constant, gives a **measure of the total number of walks** between nodes *i* and *j*

-Matrix Functions and Walks

Counting Walks: matrix functions

- ^{\circ} For **suitable choices** of $\{c_k\}_{k\geq 0}$ the series overpage converges
 - for example, $c_k = \{\delta^k\}_{k \ge 0}$, gives the matrix resolvent $(I \delta A)^{-1}$
 - different choices of *c_k* allow for different scalings

Counting Walks: matrix functions

- ^{\circ} For **suitable choices** of $\{c_k\}_{k\geq 0}$ the series overpage converges
 - for example, $c_k = \{\delta^k\}_{k \ge 0}$, gives the matrix resolvent $(I \delta A)^{-1}$
 - different choices of *c_k* allow for different scalings

In particular, if f is defined on the spectrum of A, we can define

$$F = f(A) = Pf(D)P^{-1}$$

(here $A = PDP^{-1}$ is the eigendecomposition)

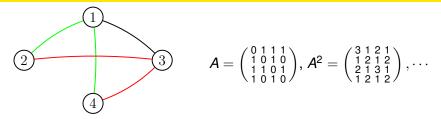
Proof Let

$$f(A) = c_0 I + \sum_{k=1}^{\infty} c_k A^k$$

substituting $A = PDP^{-1}$ into the above

$$f(A) = f(PDP^{-1}) = c_0 I + \sum_{k=1}^{\infty} c_k (PDP^{-1})^k = P\left(c_0 I + \sum_{k=1}^{\infty} c_k D^k\right) P^{-1}$$

Example: communicability

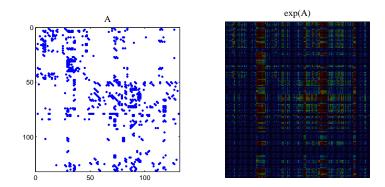


Communicability between distinct nodes *i* and *j*

$$\left(I + A + A^2/2! + A^3/3! + \cdots\right)_{ij}$$

that is $(e^A)_{ij}$ Spectral form: $(\sum_{k=1}^{n} e^{\lambda_k} \mathbf{x}^{[k]}(i) \mathbf{x}^{[k]}(j))_{ij}$

Example: communicability



Communicability applied to the frontal network of C. elegans

J. J. Crofts ()

- -Matrix Functions and Walks
- Approximate Bipartite Substructures

Outline

- 1 Introduction and Motivation
- 2 Algebraic Graph Theory
- 3 Network Reorderings
 - Directed Hierarchies
- 4 Matrix Functions and Walks
 - Approximate Bipartite Substructures
 - Weighted Networks
 - Strokes Vs Controls
- 5 Connecting it all Together

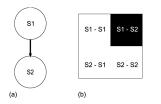
- Approximate Bipartite Substructures

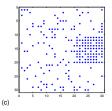
Approximate Directed Bipartite Community

Crofts, Estrada, Higham, Taylor Elec. Trans. Numer. Anal (2010)

Distinct subsets of nodes S_1 and S_2 such that

- S₁ has few internal links
- S₂ has few internal links
- there are many $S_1 \rightarrow S_2$ links
- few other links involve S₁ or S₂





-Matrix Functions and Walks

- Approximate Bipartite Substructures

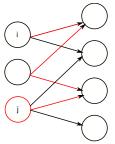
An alternating walk of length k from node i_1 to node i_{k+1}

is a list of nodes

$$i_1, i_2, i_3, \ldots, i_{k+1}$$

such that $a_{i_s,i_{s+1}} \neq 0$ for s odd, and $a_{i_{s+1},i_s} \neq 0$ for s even

Loosely, an alternating walk is a traversal that successively follows links in the forward and reverse directions



Matrix Functions and Walks

- Approximate Bipartite Substructures

This Motivates ...

$$f(A) = I - A + \frac{AA^{T}}{2!} - \frac{AA^{T}A}{3!} + \frac{AA^{T}AA^{T}}{4!} - \cdots$$

Overall idea: $f(A) + f(A^T)$ has

- positive values representing inter-community $S_1 \leftrightarrow S_1$ and $S_2 \leftrightarrow S_2$ relationships, and
- negative values representing extra-community $S_1 \leftrightarrow S_2$ relationships

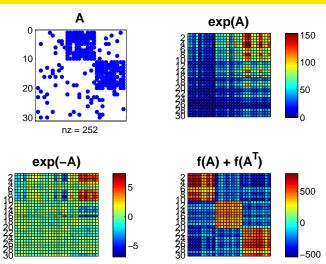
Also, $f(A) + f(A^T)$ is a symmetric matrix, so amenable to standard clustering techniques

Note: f(A) defined above is **not** a matrix function

-Matrix Functions and Walks

Approximate Bipartite Substructures

Synthetic Example

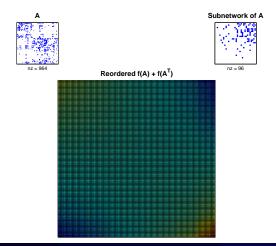


-Matrix Functions and Walks

- Approximate Bipartite Substructures

C. elegans Neural Data

Automates the computations of Durbin, (PhD thesis, Cambridge, 1987)



-Matrix Functions and Walks

-Weighted Networks

Outline

- 1 Introduction and Motivation
- 2 Algebraic Graph Theory
- 3 Network Reorderings
 - Directed Hierarchies

4 Matrix Functions and Walks

- Approximate Bipartite Substructures
- Weighted Networks
 - Strokes Vs Controls

5 Connecting it all Together

-Matrix Functions and Walks

-Weighted Networks

Communicability for a Weighted Network

W is symmetric with non-negative real weights Let

$$d_i = \sum_{k=1}^n w_{ik} \quad \& \quad D := \operatorname{diag}(d_i)$$

Normalisation:

$$W \mapsto D^{-1/2} W D^{-1/2}$$

for at least 2 reasons

- avoid overflow
- to stifle promiscuous nodes

-Matrix Functions and Walks

-Weighted Networks

Communicability for a Weighted Network

W is symmetric with non-negative real weights Let

$$d_i = \sum_{k=1}^n w_{ik} \quad \& \quad D := \operatorname{diag}(d_i)$$

Normalisation:

$$W \mapsto D^{-1/2} W D^{-1/2}$$

for at least 2 reasons

- avoid overflow
- to stifle promiscuous nodes

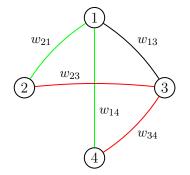
Communicability measure:

$$\exp\left(D^{-1/2}WD^{-1/2}
ight)$$

-Matrix Functions and Walks

-Weighted Networks

But What Does it Actually Mean?



W =	(0	W ₁₂	W ₁₃	W ₁₄	
	W ₂₁	0	W ₂₃	W ₂₄	
	W ₃₁	W ₃₂	0	W ₃₄	
	\ w ₄₁	W ₄₂	W43	0	Ϊ

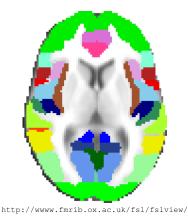
The *k*th powers of *W* provide a **measure of the total strength** contained within walks of length *k* between nodes *i* and *j*:

$$(W^2)_{ij} = \sum_k w_{ik} w_{kj}, (W^3)_{ij} = \cdots$$

- Weighted Networks

Example: anatomical connectivity data

- 9 subjects at least 6 months following first, left-hemisphere, subcortical stroke; and 10 (18) age matched controls
- Diffusion Tensor Imaging computes all connections between all voxels
- Connectivity network based on the Harvard-Oxford cortical and subcortical structural atlas: 48 cortical regions and 8 subcortical regions

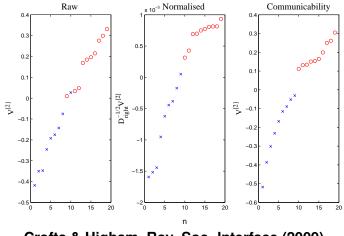


atlas-descriptions.html

-Matrix Functions and Walks

-Weighted Networks

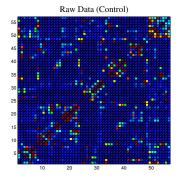
Unsupervised Clustering of Patients

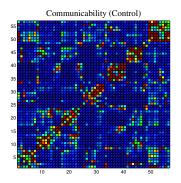


Crofts & Higham, Roy. Soc. Interface (2009)

Weighted Networks

Communicability Adds Value to the Raw Data





-Matrix Functions and Walks

Weighted Networks

What Are we Actually Detecting?

Question

Exactly what enables us to differentiate between strokes & controls?

-Matrix Functions and Walks

-Weighted Networks

What Are we Actually Detecting?

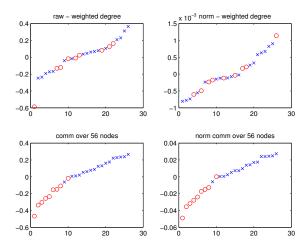
Question

Exactly what enables us to differentiate between strokes & controls?

- If it is merely the fact that many connections have been destroyed close to the infarcted region this is not very interesting – an MRI scan can tell us this with no further analysis needed!
 - Why not rerun the tractographies including only those connections within the non-stroke hemisphere?
 - If we still distinguish between the two classes, this will be more relevant from a bio point of view

-Weighted Networks

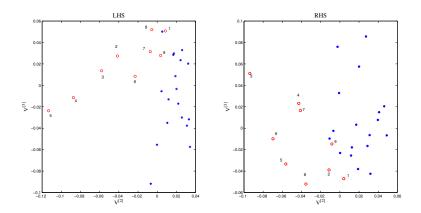
RHS Hemisphere Sorted by Brain Region



-Matrix Functions and Walks

-Weighted Networks

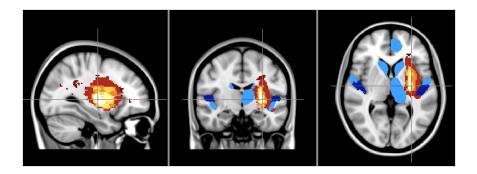
Left (Stroke Side) Vs Right



-Matrix Functions and Walks

-Weighted Networks

Overlay of Stroke Overlap Volume

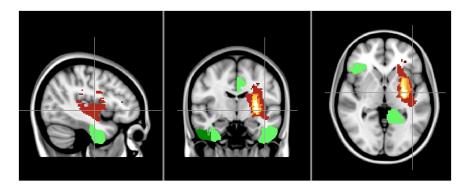


Highlighted regions are those where patients showed reduced communicability relative to controls

-Matrix Functions and Walks

Weighted Networks

Crofts et al. NeuroImage (2011)



And increased communicability

J. J. Crofts ()

UCSB 2011

-Matrix Functions and Walks

-Weighted Networks



[©] Generalised singular value decomposition

$$A = UCX^{-1}$$
 and $B = VSX^{-1}$

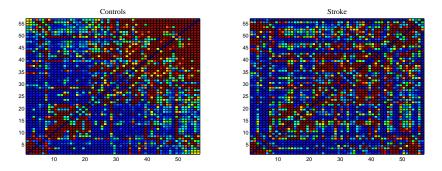
of a pair of matrices can be used to determine clusters that are 'good' in one network and 'poor' in the other - and vice versa **[Xiao et al. (2011)]**

Another way to understand this is to note that, in the case where A and B are invertible the GSVD is closely related to the SVD of AB⁻¹ and BA⁻¹, hence

$$AB^{-1} = UCS^{-1}V^T$$
 and $BA^{-1} = VSC^{-1}U^T$

- Weighted Networks

GSVD Example: anatomical connectivity data



GSVD finds a group of brain regions that are much better connected in controls than strokes

-Connecting it all Together

Outline

- Introduction and Motivation
- 2 Algebraic Graph Theory
- 3 Network Reorderings
 - Directed Hierarchies
- 4 Matrix Functions and Walks
 - Approximate Bipartite Substructures
 - Weighted Networks
 - Strokes Vs Controls
- 5 Connecting it all Together

- Connecting it all Together

The Take Home Message

- Spectral methods provide a useful tool for determining different types of network architecture
 - scalable
 - most work has focused on the adjacency matrix

- Connecting it all Together

The Take Home Message

- Spectral methods provide a useful tool for determining different types of network architecture
 - scalable
 - most work has focused on the adjacency matrix
- Walk based measures are attractive since:
 - combinatorics are described in terms of basic linear algebra
 - information does not necessarily flow along geodesics
 - walks are more tolerent to errors

- Connecting it all Together

The Take Home Message

- O Spectral methods provide a useful tool for determining different types of network architecture
 - scalable
 - most work has focused on the adjacency matrix
- Walk based measures are attractive since:
 - combinatorics are described in terms of basic linear algebra
 - information does not necessarily flow along geodesics
 - walks are more tolerent to errors

[©] GSVD and extensions:

- tensor decompositions (plus multiple tasks/modalities, time-dependent networks)
- NNMF (Lee et al (2010)), ICA (Smith et al (2005))

- Connecting it all Together

Thank you!

Colleagues at Strathclyde: Des Higham, Ernesto Estrada & Alan Taylor

Colleagues from Oxford: Heidi Johanson-Berg & Tim Behrens The stroke data was supplied by Rose Bosnell

This work was supported by the **Medical Research Council** under project no. MRC G0601353